## Minimizing State Preparations for VQE

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## Results



## Background: Ground State Estimation

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- Classically, diagonalize exponentially-sized matrix.
- Quantum Phase Estimation algorithm showed how to solve in poly-time.


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- But, commuting terms can be measured simultaneously.


## Our Contributions

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- How to group terms into large commuting families?
- Impact on variance and study of covariances
- Benchmarking \& resource estimation for representative molecules


## Pauli Commutativity Relations

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## Commutator Notation

$[A, B]=A B-B A \begin{cases}=0 & \text { if } A \text { and } B \text { commute } \\ \neq 0 & \text { if } A \text { and } B \text { do not commute }\end{cases}$

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- Other Pauli commutators follow from cyclic multiplication property.:
- $X Y=i Z=-Y X$
- $Y Z=i X=-Z Y$
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## Examples

$$
\text { QWC: } \begin{gathered}
X \\
\text { I Y Z Z Y I I Y } \\
\text { I }
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\begin{gathered}
\text { QWC: } \left.\begin{array}{cccccc}
X & Y & Z & Y & I & Y \\
I & Y & Z & I & I & I
\end{array} \right\rvert\, \\
\text { Not QWC: } \\
\begin{array}{llllll}
X & Y & Z & Y & I & Y \\
I & Y & Z & X & I & I
\end{array}
\end{gathered}
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| :--- | :--- | :--- |
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| I | $X$ | $Z$ |
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These terms are a QWC family. Measure in:

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\begin{array}{lll}
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\text { I } & X & I \\
\text { I } & X & Z \\
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Z & I & Z \\
Z & X & I \\
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\end{array}
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These terms are a QWC family. Measure in:

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In general: QWC simultaneous measurements requires $O(N)$ single qubit gates (depth $=1$ ).

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## General Commutativity of Pauli Strings

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## General Commutativity of Pauli Strings

If an even number of indices don't commute, then $A$ and $B$ commute. Proof: All Pauli strings are either commuting or anti-commuting, so we know that either $A B=B A$ or $A B=-B A$. Each commuting index multiplies by +1 , each non-commuting index multiplies by -1 . Need even number of non-commuting indices to have a total +1 .

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[Devoret Notes]

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(3) Apply $O\left(N^{2}\right)$ gates from resulting circuit (our software gives explicit circuit decomposition):
(9) Result: Bell Basis Measurement for this example
$O\left(N^{2}\right)$ gates is fine, because UCCSD ansatz prep is $O\left(N^{3}\right)$ or $O\left(N^{4}\right)$.

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We seek MIN-CLIQUE-COVER.
Problem is NP-Hard, but use heuristics (Boppana-Halldórsson). Note that our problem is also NP-HARD by reduction from MIN-CLIQUE-COVER.

## Results: Across Molecules



## Results: Across Encodings

Term Grouping for H2 (6-31g basis, 8 modes)


## Results: Across Active Spaces



## Qubit Tapering

- Create QWCommutativity by transforming Hamiltonian [Bravyi 2017].


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| $\sigma_{1}^{z}$ | $\sigma_{2}^{z}$ | $\sigma_{3}^{z}$ | $\sigma_{4}^{z}$ |
| :--- | :--- | :--- | :--- |
| $\sigma_{1}^{z} \sigma_{2}^{z}$ | $\sigma_{1}^{z} \sigma_{3}^{z}$ | $\sigma_{1}^{z} \sigma_{4}^{z}$ |  |
| $\sigma_{2}^{z} \sigma_{3}^{z}$ | $\sigma_{2}^{z} \sigma_{4}^{z}$ | $\sigma_{3}^{z} \sigma_{4}^{z}$ |  |
| $\sigma_{1}^{y} \sigma_{2}^{y} \sigma_{3}^{x} \sigma_{4}^{x}$ | $\sigma_{1}^{x} \sigma_{2}^{y} \sigma_{3}^{y} \sigma_{4}^{x}$ | $\sigma_{1}^{y} \sigma_{2}^{x} \sigma_{3}^{x} \sigma_{4}^{y}$ | $\sigma_{1}^{x} \sigma_{2}^{x} \sigma_{3}^{y} \sigma_{4}^{y}$ |

$$
\begin{aligned}
& U_{1}=\frac{1}{\sqrt{2}}\left(\sigma_{2}^{x}+\sigma_{1}^{z} \sigma_{2}^{z}\right), \quad U_{2}=\frac{1}{\sqrt{2}}\left(\sigma_{3}^{x}+\sigma_{1}^{z} \sigma_{3}^{z}\right) \\
& \text { and } U_{3}=\frac{1}{\sqrt{2}}\left(\sigma_{4}^{x}+\sigma_{1}^{z} \sigma_{4}^{z}\right)
\end{aligned}
$$

| $\sigma_{1}^{z}$ | $\sigma_{1}^{z} \sigma_{2}^{x}$ | $\sigma_{1}^{z} \sigma_{3}^{x}$ | $\sigma_{1}^{z} \sigma_{4}^{x}$ |
| :--- | :--- | :--- | :--- |
| $\sigma_{2}^{x}$ | $\sigma_{3}^{x}$ | $\sigma_{4}^{x}$ |  |
| $\sigma_{2}^{x} \sigma_{3}^{x}$ | $\sigma_{2}^{x} \sigma_{4}^{x}$ | $\sigma_{3}^{x} \sigma_{4}^{x}$ |  |
| $\sigma_{1}^{x} \sigma_{3}^{x} \sigma_{4}^{x}$ | $\sigma_{1}^{x} \sigma_{4}^{x}$ | $\sigma_{1}^{x} \sigma_{2}^{x} \sigma_{3}^{x}$ | $\sigma_{1}^{x} \sigma_{2}^{x}$ |

- In example, three qubits are tapered out of H2 Hamiltonian


## Qubit Tapering Results

## Hamiltonian \# of Qubits \# Tapered <br> H2 (small \# active spaces) 4 <br> 3 <br> H2 (large \# active spaces) $8 \quad 2$ <br> H 2 O <br> 8 <br> 2

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${ }^{1}$ Example from [McClean et al 2015]

## $k=5$ Groups



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$$
\begin{array}{r}
E(\# \text { state preps })= \\
k(\operatorname{Var}(-X X)+\operatorname{Var}(-Y Y)+\operatorname{Var}(Z Z)+\operatorname{Var}(Z I)+\operatorname{Var}(I Z)) / \epsilon^{2} \\
=5(1+1+0+0+0) / \epsilon^{2} \\
=10 / \epsilon^{2}
\end{array}
$$

## $k=3$ Groups



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$$
\begin{array}{r}
E(\# \text { state preps }) / \epsilon^{2}= \\
k[\operatorname{Var}(-X X)+\operatorname{Var}(\{-Y Y,-Z Z\})+\operatorname{Var}(\{Z I, I Z\})] / \epsilon^{2} \\
=k[\operatorname{Var}(-X X)+(\operatorname{Var}(-Y Y)+\operatorname{Var}(-Z Z)+2 \operatorname{Cov}(-Y Y,-Z Z)) \\
+(\operatorname{Var}(Z I)+\operatorname{Var}(I Z)+2 \operatorname{Cov}(I Z, Z I))] / \epsilon^{2} \\
=3[1+(1+0+0)+(0+0+0)] / \epsilon^{2}=6 / \epsilon^{2}
\end{array}
$$

## $k=2$ Groups



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$$
\begin{array}{r}
E(\# \text { state preps })= \\
k[\operatorname{Var}(\{-X X,-Y Y, Z Z\})+\operatorname{Var}(\{Z I, I Z\})] / \epsilon^{2} \\
=k[(\operatorname{Var}(-X X)+\operatorname{Var}(-Y Y)+\operatorname{Var}(Z Z)+ \\
2 \operatorname{Cov}(-X X,-Y Y)+2 \operatorname{Cov}(-X X, Z Z)+2 \operatorname{Cov}(-Y Y, Z Z)) \\
(\operatorname{Var}(Z I)+\operatorname{Var}(I Z)+2 \operatorname{Cov}(I Z, Z I))] / \epsilon^{2} \\
=2[(1+1+0+2 * 1+0+0)+(0+0+0)] / \epsilon^{2}=8 / \epsilon^{2}
\end{array}
$$

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Work in progress: adaptively adjust groups after initial grouping.

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## Thanks!


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